COSC/MATH 4P61 - Theory of Computation Example Questions Test 2

Question 1: Consider the DFA below. Construct the minimum-state equivalent DFA.



Question 2: Use the Pumping Lemma to show that the language

$$L = \{0^i 1^{2i} \mid i \ge 1\}$$

is not regular.

Question 3: Suppose the following PDA $P = (\{q, r\}, \{0, 1\}, \{Z_0, X\}, \delta, q, Z_0, \{r\})$

is given:



Convert P to a PDA P' with N(P') = L(P).

Question 4: Suppose the following PDA $P = (\{q, r\}, \{0, 1\}, \{Z_0, X\}, \delta, q, Z_0, \emptyset)$ is given:



Convert P to a PDA P' with L(P') = N(P).

Question 5: Convert the grammar

$$\begin{split} S &\to 0S1 \mid A \\ A &\to 1A0 \mid S \mid \epsilon \end{split}$$

to a PDA that accepts the same language by empty stack.

Question 6: Consider the following grammar:

$$\begin{split} S &\to ASB \mid \epsilon \\ A &\to aAS \mid a \\ B &\to SbS \mid A \mid bb \end{split}$$

- a. Eliminate all ϵ -productions.
- b. Eliminate all unit productions from the resulting grammar in a).
- c. Eliminate all useless symbols from the resulting grammar in b).
- d. Put the resulting grammar in c) in Chomsky Normal Form.

Solutions

Question 1:



Question 2: Assume *L* is regular and let *n* be the constant from the Pumping Lemma. Pick the word $w = 0^{n}1^{2n}$. Using the Pumping Lemma we can write w as w = xyz with $|xy| \le n$. Therefore $y \ne \epsilon$ consists of only 0's. The word $xy^{0}z = xz$ must be in *L*, which is a contradiction since it contains 2n copies of 1 but less than *n* copies of 0.

Question 3: The PDA $P' = (\{p_0, p, q, r\}, \{0, 1\}, \{X_0, Z_0, X\}, \delta, p_0, X_0, \emptyset)$ is given by:



Question 4: The PDA $P' = (\{p_0, p_f, q, r\}, \{0, 1\}, \{X_0, Z_0, X\}, \delta, p_0, X_0, \{p_f\})$ is given by:



Question 5: The PDA $P = (\{q\}, \{0, 1\}, \{0, 1, S, A\}, \delta, q, S, \emptyset)$ is given by:



Question 6:

a. Only S is nullable so that we obtain:

$$S \rightarrow ASB \mid AB$$

$$A \rightarrow aAS \mid aA \mid a$$

$$B \rightarrow SbS \mid bS \mid Sb \mid b \mid A \mid bb$$

b. There is only one unit production $B \to A$ so that we obtain:

$$S \rightarrow ASB \mid AB$$

$$A \rightarrow aAS \mid aA \mid a$$

$$B \rightarrow SbS \mid bS \mid Sb \mid b \mid aAS \mid aA \mid a \mid bb$$

- c. A and B are obviously generating. From $S \to AB$ we conclude that also S is generating. There are no unreachable symbols. As a consequence all symbols are useful.
- d. We obtain the following grammar in Chomsky Normal Form:

$$\begin{array}{l} S \rightarrow AE \mid AB \\ A \rightarrow CF \mid CA \mid a \\ B \rightarrow SG \mid DS \mid SD \mid b \mid CF \mid CA \mid a \mid DD \\ C \rightarrow a \\ D \rightarrow b \\ E \rightarrow SB \\ F \rightarrow AS \\ G \rightarrow DS \end{array}$$