### PARALLEL PROGRAMMING RECIPE

The process of development of parallel programs can be summarized as follows:

- 1. Pick up a particular problem of interest;
- 2. Conceptualize the solution;
- 3. Split this solution into components to be executed simultaneously as cooperating processes;
- 4. Code each component;
- 5. Arrange components in groups;
- 6. Allocate to each group a separate processor of suitable type;
- 7. Execute simultaneously all components, noting overall run time.

A question might be asked:

WHY GROUPS?

Another nasty question:

We hope for speedup > 1. What if the application of this recipe yields speedup < 1 ? In particular, speedup = 0 implies deadlock!

# WHY PROCESS GROUPS?

We may face two possible situations:

 If the number of processes <= number of processors and it is possible to satisfy all processes' needs for processor types,

#### then

Allocate each process to a separate processor, Keeping all processor types compatible with processes' needs

### 2. else

We must have groups of processes, Each group being allocated a single processor, and All processes within each group running in a time-sharing mode.

NOTE 1: We focus on situation (2), considering situation (1) as special case of situation (2).

NOTE 2: In general, no brute-force processor allocation approach is feasible, because there are  $N^P$  ways to allocate P processes among N processors.

However, this approach is feasible for a small N and P, and the evolution seems to "know" it...

# THE PROCESSOR ALLOCATION PROBLEM

### Given:

1. A computer with N processors (of possibly differing characteristics like speed, amount of local memory, presence or absence of floating point facilities, graphics, FFT, etc.)

### and

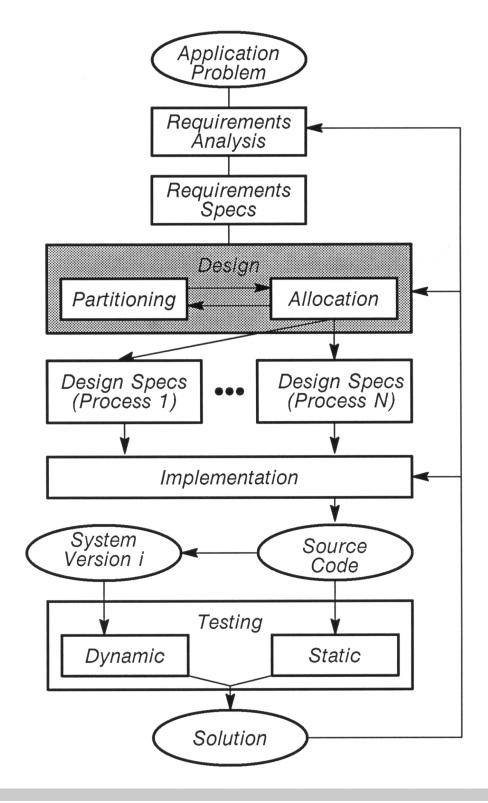
2. An algorithm consisting of a mix of P processes (of possibly differing processor type preferences)

### Decide:

- 1. Which process(es) to run on which processor(s), when, and for how long?
- 2. On the order sequence of such assignments;
- 3. On the computer architecture, facilitating interprocess communication.

Our measure of success will be calculated according to some selected criterion (like speedup, node efficiency, area efficiency, etc.)

# DISTRIBUTED SOFTWARE LIFE CYCLE



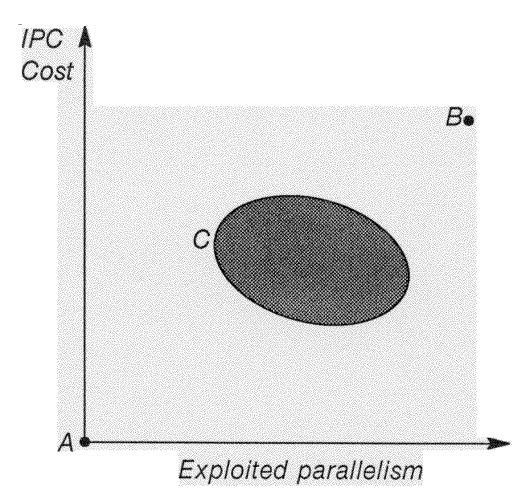
### PARTITIONING CONSIDERATIONS

**OBJECTIVES**:

- Minimize interprocess communication
- Exploit potential concurrency
- Limit sizes of processes

### DIFFICULTIES:

- · How to measure effectiveness before allocation?
- Partitioning criteria conflicts.



# TASK ALLOCATION CONSIDERATIONS

Combinatorially, if we have P processes and N processors, there are  $N^P$  possible allocations (or more, if we allow replication to enhance fault tolerance)

Allocation effectiveness depends on the allocation GOAL, like:

- Minimize total IPC cost
- Minimize total computation and IPC cost
- Minimize completion time
- Minimize load imbalance
- Maximize system reliability

Feasible allocations must meet system constraints, like:

- Memory capacity
- Processing time limits